

# Some remarks on the asymptotic behavior of cyclic AG-codes

R. Toledano<sup>1</sup>

Facultad de Ingeniería Química - Univ. Nacional del Litoral

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<sup>1</sup>Joint work with M. Chara and R. Podestá

# Linear codes

- A **linear code**  $\mathcal{C}$  (over the alphabet  $\mathbb{F}_q$ ) of **length**  $n$  is a linear subspace of  $\mathbb{F}_q^n$ . The elements of  $\mathcal{C}$  are usually called **codewords**.
- If  $k = \dim \mathcal{C}$  is the **dimension** of the code (as a vector space over  $\mathbb{F}_q$ ) and

$$d = \min\{d(c, c') : c, c' \in \mathcal{C}, c \neq c'\}$$

is the **minimum distance** of  $\mathcal{C}$ , where  $d$  is the Hamming distance in  $\mathbb{F}_q^n$ , we shall say that  $\mathcal{C}$  is an  $[n, k, d]$ -code.

- There is a natural action of the permutation group  $\mathbb{S}_n$  on  $\mathbb{F}_q^n$  given by

$$\pi \cdot (v_1, \dots, v_n) = (v_{\pi(1)}, \dots, v_{\pi(n)})$$

where  $\pi \in \mathbb{S}_n$  and  $(v_1, \dots, v_n) \in \mathbb{F}_q^n$ .

- The set of all  $\pi \in \mathbb{S}_n$  such that  $\pi \cdot c \in \mathcal{C}$  for all codewords  $c$  of  $\mathcal{C}$  forms a subgroup  $\text{Perm}(\mathcal{C})$  of  $\mathbb{S}_n$  which is called the **permutation group of  $\mathcal{C}$** :

$$\text{Perm}(\mathcal{C}) = \{\pi \in \mathbb{S}_n : \pi(\mathcal{C}) = \mathcal{C}\}.$$

# Transitive and cyclic codes

Let  $c = (c_1, \dots, c_{n-1}, c_n)$  a codeword in  $\mathcal{C}$ .

- $\mathcal{C}$  is **transitive** if  $\text{Perm}(\mathcal{C})$  acts transitively on  $C$ , i.e. if for any  $1 \leq i < j \leq n$  there is a  $\pi \in \mathbb{S}_n$  such that  $\pi(i) = j$ .
- $\mathcal{C}$  is **cyclic** if it is invariant under the action of the cyclic shift  $s \in \mathbb{S}_n$  defined as  $s(1) = n$  and  $s(i) = i - 1$  for  $i = 2, \dots, n$ , i.e. if

$$s \cdot c = (c_n, c_1, \dots, c_{n-1}) \in \mathcal{C}$$

for every  $c \in \mathcal{C}$ .

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**Cyclic codes are transitive codes**

## AG-codes

- Let  $F$  be an algebraic function field over  $\mathbb{F}_q$ , i.e. a finite extension of a rational function field  $\mathbb{F}_q(x)$ .
- Let  $D$  and  $G$  be disjoint divisors of  $F$ , with  $D = P_1 + \dots + P_n$  where  $P_1, \dots, P_n$  are different rational places.
- Let  $\mathcal{L}(G)$  be the **Riemann-Roch space associated to  $G$**

$$\mathcal{L}(G) = \{x \in F : (x) \geq -G\} \cup \{0\}$$

where  $(x)$  denotes the principal divisor associated to  $x \in F$ .

**The AG-code defined by  $F$ ,  $D$  and  $G$  is**

$$\mathcal{C}(D, G) = \{(x(P_1), \dots, x(P_n)) \in \mathbb{F}_q^n : x \in \mathcal{L}(G)\}$$

**where  $x(P_i)$  stands for the residue class of  $x$  modulo  $P_i$ .**

## Cyclic AG-codes

An AG-code  $C_{\mathcal{L}}(D, G)$  with  $D = P_1 + \cdots + P_n$  is cyclic if for any codeword

$$(x(P_1), \dots, x(P_n)) \in C_{\mathcal{L}}(D, G),$$

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This happens if and only if there exists  $z \in \mathcal{L}(G)$  such that

$$\begin{aligned} z(P_1) &= x(P_n), \\ z(P_2) &= x(P_1), \\ &\vdots \\ z(P_n) &= x(P_{n-1}). \end{aligned} \tag{1}$$



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The existence of such an element  $z \in \mathcal{L}(G)$  is a crucial question to answer in the theory of cyclic AG-codes.

Let  $E/F$  a finite cyclic extension of the function field  $\mathbb{F}_q$  and  $\mathcal{C}$  an AG-code such that

- every place in  $\text{Sup}(D) = \{P_1, \dots, P_n\} \subset \mathbb{P}(E)$  lies over a unique place  $P$  of  $F$ ;
- there is an element  $\sigma \in \text{Aut}(E/\mathbb{F}_q(x))$  such that

$$\sigma(G) = G \quad \text{and} \quad \sigma(P_i) = P_{i-1 \bmod n},$$

for  $i = 1, \dots, n$ .

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**$\mathcal{C}$  is a cyclic code**

Take  $z = \sigma^{-1}(x)$ .

# Cyclic AG-codes and cyclic extension

The previous situation can happen only in the presence of cyclic extensions:

## Proposition

*Let  $F$  be a function field over  $\mathbb{F}_q$  and let  $P_1, \dots, P_n$  be  $n$  different places of  $F$ . Suppose there exists  $\sigma \in \text{Aut}(F/\mathbb{F}_q(x))$  such that  $\sigma(P_1) = P_n, \sigma(P_2) = P_1, \dots, \sigma(P_n) = P_{n-1}$ . Then there exist an intermediate field  $\mathbb{F}_q(x) \subset E \subset F$  and a place  $P$  of  $E$  such that  $F/E$  is a cyclic extension of degree  $m$  divisible by  $n$ , and  $P$  decomposes exactly in  $F$  into the places  $P_1, \dots, P_n$  with  $e(P_i|P)f(P_i|P) = \frac{m}{n}$  for  $i = 1, \dots, n$ .*

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*Conversely, let  $F/E$  be a cyclic extension of function fields over  $\mathbb{F}_q$  of degree  $m$ . Let  $P$  be a place of  $E$  and let  $P_1, \dots, P_n$  be all the places of  $F$  lying over  $P$ . Then,  $m$  is divisible by  $n$  and we have that  $\sigma(P_1) = P_n, \sigma(P_2) = P_1, \dots, \sigma(P_n) = P_{n-1}$  for any generator  $\sigma$  of  $\text{Gal}(F/E)$ .*

# Towers of function fields

A sequence  $\mathcal{F} = \{F_i\}_{i=0}^{\infty}$  of function fields over  $\mathbb{F}_q$  is called a **tower** if

- $F_0 \subsetneq F_1 \subsetneq \cdots \subsetneq F_i \subsetneq \cdots$ ;
- $F_{i+1}/F_i$  finite and separable,  $i \geq 0$ ;
- $\mathbb{F}_q$  is the full constant field of  $F_i$ ,  $i \geq 0$ .
- $g(F_i) \rightarrow \infty$  for  $i \rightarrow \infty$

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$$\lambda(\mathcal{F}) := \lim_{i \rightarrow \infty} \frac{N(F_i)}{g(F_i)} \geq 0.$$



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**The tower is asymptotically good if  $\lambda(\mathcal{F}) > 0$ .**

## Theorem

Let  $\mathcal{F} = \{F_i\}_{i=0}^{\infty}$  be an asymptotically good tower of function fields over  $\mathbb{F}_q$  where  $F_0 = \mathbb{F}_q(x)$  is a rational function field. For each  $i \in \mathbb{N}$ , let  $n_i$  be a positive integer and let  $P_1, \dots, P_{n_i}$  be different rational places of  $F_i$ . Suppose that for each  $i \in \mathbb{N}$  there is an element  $\sigma_i \in \text{Aut}(F_i/F_0)$  such that

$$\sigma(P_1) = P_{n_i}, \sigma(P_2) = P_1, \dots, \sigma(P_{n_i}) = P_{n_i-1}.$$

Then  $n_i < [F_i : F_0]$  and there exists a place  $P \in \text{Ram}(\mathcal{F})$  such that  $P_1, \dots, P_{n_i}$  are all the places of  $F_i$  lying over  $P$ .

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This theorem shows that the divisor  $D$  has to be defined with all the rational places lying over a place in the ramification locus of the tower.

**Thanks!**